

Machine Problem 3 – CPU scheduling

Deadline:2016/12/18 23:59

● Problem description:

The default CPU scheduling algorithm of Nachos is a simple round-robin scheduler for every 500 ticks. The goal of MP3 is to replace it with a multilevel feedback queue as described below.

- There are **3 levels of queues**: L1, L2 and L3. L1 is the highest level queue, and L3 is the lowest level queue. The next scheduling job is always selected from the highest level queue with available jobs.
- All processes must have a valid **scheduling priority between 0 to 149**. Higher value means higher priority. So 149 is the highest priority, and 0 is the lowest priority.
- A process with priority between **0~49 is in L3** queue. A process with priority between **50~99 is in L2** queue. A process with priority between **100~149 is in L1** queue.
- **L1 queue uses a SJF**(shortest job first) scheduling algorithm. The job execution time is approximated using the equation: $t(i) = 0.5 \cdot T + 0.5 \cdot t(i - 1)$
- **L2 queue uses a priority** scheduling algorithm.
- **L3 queue uses a round-robin** scheduling algorithm with time quantum **100 ticks** instead of 500 ticks.
- An **aging mechanism** must be implemented, so that the priority of a process is **increased by 10 after waiting for every 1500 ticks**.
- Allow the **initial priority** of a process to be set by reading the input argument“-ep” from Nachos command line.

E.g.,: the command below will launch 2 processes: test2 with initial priority 40, and test3 with initial priority 80.

“../build.linux/nachos -ep test2 40 -ep test3 80”

- Working items:

1. (30%) L1 SJF scheduling algorithm as described above.
2. (15%) L2 priority job scheduling algorithm as described above.
3. (5%) L3 round-robin scheduling algorithm as described above.
4. (15%) An aging mechanism to move processes among the queues as described above, and a “-ep” input argument for setting initial priority.
5. (5%) Output log information during your execution: **(you must follow the exact output format as below)**
 - (1). Whenever a process is insert into a queue:
Tick [current tick count]: Thread [thread ID] is inserted into queue L[queue level]
 - (2). Whenever a process is removed from a queue:
Tick [current tick count]: Thread [thread ID] is removed from queue L[queue level]
 - (3). Whenever a process changes its scheduling priority:
Tick [current tick count]: Thread [thread ID] changes its priority from [old value] to [new value]
 - (4). Whenever a context switch occurs
Tick [current tick count]: Thread [new thread ID] is now selected for execution
Tick [current tick count]: Thread [prev thread ID] is replaced, and it has executed [tick count] ticks
6. (15%) Report to explain your implementation and basic team info.
7. (15%) Demo.
 - (1). You must prepare test cases yourself to demonstrate the correctness of your implementation from item1 to item4. One test case per working item is preferred.
 - (2). The correctness of your demonstration must be proven ONLY using the output log information from working item5.
 - (3). Random test case will also be given during the demo to verify the correctness.

- Hint:

The following files “may” be modified...

1. threads/ main.*, kernel.*, threads.*, scheduler.*, alarm.*
2. machine/ interrupt.*, timer.*, stats.*